Logic R, list 3. (homework due Nov 29, declarations due by 13:00 on Nov 29)

- 1. Let  $BA_0$  be the theory of atomless Boolean algebras. Prove that  $BA_0$  is  $\aleph_0$ -categorical and admits elimination of quantifiers.
- 2. Prove that if T is complete and admits elimination of quantifiers, then for  $M, N \models T, M \subseteq N$  implies  $M \prec N$ .
- 3. Check that the Stone space S(B) of Boolean algebra B is compact Hausdorff. Prove that if B is countable, then S(B) is homeomorphic to a closed subset of the Cantor set.
- 4. Assume that M is a model (structure) for language L and N is a model of the theory Th(M, n)<sub>m∈M</sub> (in language L(M), with new constant symbols naming all elements of |M| added). Let M' = {m<sup>N</sup> : m ∈ M}. Prove that
  (a) M' is a model for L (with the structure induced from N).
  (b) M' ≅ M and M' ≺ N.
- 5. Prove the every proper filter in a Boolean algebra B extends to an ultrafilter.
- 6. (a) Prove that every two Boolean algebras of the same size are isomorphic.
  (b) Prove that a Boolean algebra B is finite ↔ S(B) is finite.
  (c) Prove that the number of elements of a finite Boolean algebra is a power of 2.
- 7. Prove that in some elementary extension  $\mathbb{Z}'$  of the group  $(\mathbb{Z}, +)$  there is an element  $a \neq 0$  that is divisible (i.e. for every n > 0,  $\mathbb{Z}' \models \exists x \underbrace{x + \cdots + x}_{n} = a$ ).
- 8. (a) Describe all countable models of the theory of (Z, S), where S is the successor function (up to isomorphism).
  (b) A cycle in a k element set X = {x, ..., x, } is a structure (X, S), where S

(b) A cycle in a k-element set  $X = \{x_1, \ldots, x_k\}$  is a structure (X, S), where S is a permutation of X such that  $S(x_i) = x_{i+1}$  for i < k, and  $S(x_k) = x_1$ . Let T be the theory of finite cycles in the language  $L = \{S\}$ , that is, the set of formulas of L true in all finite cycles (X, S). Prove that T and  $Th(\mathbb{Z}, S)$  have the same infinite countable models.

9. Assume that G is a group of permutations of an non-empty finite set X (i.e. G < S(X)). Prove that for certain language L, there is an L-structure on the set X such that G = Aut(X).