

ex. 2 Suppose T is complete and satisfies elimination of quantifiers, $M, N \models T$ and $M \subseteq N$, then $M \models N$.

Proof

First, let us recall the definition of a substructure.

$$M \subseteq N \stackrel{\text{def}}{\Leftrightarrow} 1) |M| \subseteq |N|$$

$$2) \text{For } n\text{-ary } f \in \lambda: f^M = f^N \upharpoonright_{M^n}$$

$$3) \text{For } n\text{-ary } R \in \lambda: R^M \subseteq R^N \cap M^n$$

$$4) \text{For constant } c: c^M = c^N.$$

Lemma. For every q.f. sentence $\varphi \in \mathcal{T}_{\lambda(M)}$ we have $M \models \varphi \Leftrightarrow N \models \varphi$ (where we think of M, N as $\lambda(M)$ -structures).

Proof of lemma Induction on complexity of φ .

First we need to say that for every constant term $t \in \mathcal{T}_{\lambda(M)}$ we have $t^M = t^N$.

This is true by the definition of
Substructure.

Induction on complexity of φ :

• $\varphi : R_i(t_1, \dots, t_n) :$

$$M \Leftrightarrow \varphi \Leftrightarrow R_i^M(t_1^M, \dots, t_n^M) \xrightarrow{\text{def. of subst.}} \quad \xleftarrow{\text{def. of subst.}}$$

$$R_i^N(t_1^N, \dots, t_n^N) \Leftrightarrow N \models \varphi$$

• $\varphi : t_1 = t_2 : M \models \varphi \Leftrightarrow t_1^M = t_2^M \Leftrightarrow t_1^N = t_2^N \Leftrightarrow N \models \varphi$

$$\begin{array}{ccc} \parallel & & \parallel \\ t_1^M & & t_2^N \\ \parallel & & \parallel \\ t_1 & & t_2 \end{array}$$

• $\varphi : \varphi_1 \wedge \varphi_2, \neg \psi : \text{trivial.}$

lemma 

Now, take any $\varphi(\bar{x}) \in \mathcal{F}_L$. Let $\psi(\bar{x}) \in \mathcal{F}_L$

be a formula without quantifiers s.t.

$$T \vdash \varphi(\bar{x}) \Leftrightarrow \psi(\bar{x}).$$

Now by our lemma we know that

for every $\bar{a} \subseteq M$:

$$M \models \varphi(\bar{a}) \Leftrightarrow M \models \psi(\bar{a}) \xrightarrow{\text{Lemma}} N \models \psi(\bar{a}) \Leftrightarrow N \models \varphi(\bar{a})$$



Ex. 3 $B - BA$.

- 1) $S(B)$ is compact Hausdorff
- 2) B b.tble $\Rightarrow S(B)$ is homeomorphic to a closed subset of the Cantor set.

First we will prove two usefull lemmes:

Lemma 1 Every subset $A \subseteq B$ with FIP
(Finite Intersection Property, i.e. $\forall \bar{\alpha} \subseteq A \wedge \bar{\alpha} \neq \emptyset \exists \alpha \in \bar{\alpha} \cap \text{finite}$)
can be extended to a proper filter.

Proof Take any subset $A \subseteq B$ with FIP.

Wlog suppose that $\forall \bar{\alpha} \subseteq A, \text{finite } (\forall \bar{\alpha} \in \bar{\alpha} \subseteq A)$.

Now let $F = A \cup \{b \in B : \exists \alpha \subseteq A (\alpha \leq b)\}$.

$F \supseteq A$ and it's a proper filter:

- It is certainly closed under \leq
- It is closed under finite intersections and $\emptyset \notin F$:

Take any finite $\bar{a} \subseteq F$. Wlog $a_1, \dots, a_k \in A$,
 $a_{k+1}, \dots, a_n \notin A$ for some k ($n = |\bar{a}|$).

By the construction of F there are

$b_{k+1}, \dots, b_n \in A$ s.t. $b_i \leq a_i$, $i = k+1, \dots, n$.

Now because A has FIP, then

$$0 \neq a_1 \wedge \dots \wedge_k \wedge b_{k+1} \wedge \dots \wedge b_n \leq \leftarrow \text{this is easy to show, 3/1 omit it.}$$

$$\leq a_1 \wedge \dots \wedge_k \wedge a_{k+1} \wedge \dots \wedge a_n.$$



Lemma 2 $F \in S(B)$, then $\forall a \in B (a \in F \vee a' \in F)$

Proof (A.a.) Suppose $F \in S(B)$ and

that there is $a \in B$ s.t. $a, a' \notin F$.

If $\forall b \in F a \wedge b \neq 0$, then $F \cup \{a\}$

has FIP, thus it could be extended

to a filter bigger than F , which

contradicts maximality of F . Similarly,

it can't be that $\forall b \in F a' \wedge b \neq 0$.

So there is $b, c \in F$ s.t. $a \wedge b = a' \wedge c = \emptyset$

But then $\emptyset = (a \wedge b) \vee (a' \wedge c) =$

$$= a' \vee b' \vee a \vee c' = b' \vee c' = (b \wedge c)'$$

↓

$$b \wedge c = \emptyset$$

↯

Now we can tackle the problems.

Proof 1) First, we'll show that $S(B)$

is Hausdorff. Take any $F_1, F_2 \in S(B)$,

$F_1 \neq F_2$. Take $a \in F_1$ s.t. $a \notin F_2$ and

$b \in F_2$ s.t. $b \notin F_1$ (there are such

a, b because $F_1 \neq F_2, F_2 \neq F_1$).

Now by lemma 2 $b' \in F_1, a' \in F_2$, so

$F_1 \in [a \wedge b']$, $F_2 \in [b \wedge a']$ and

$$[a \wedge b'] \cap [a' \wedge b] = \emptyset.$$

□

Now let's show that $S(B)$ is compact.

It's sufficient to show that every cover consisting exclusively of base

open sets has a finite subset that's also a cover of $S(B)$. Let's take any $A \subseteq B$ s.t. $M = \{[\alpha] : \alpha \in A\}$ is a cover. (A.a.) Suppose that there's no such finite subset of M that's also a cover. That is, for every finite $\bar{A} \subseteq A$ there is $F \in S(B)$ s.t. $F \not\in \bigcup_{i=1}^n [\alpha_i]$.

What it means is that $\alpha_i \notin F$ for $i = 1, \dots, n$. From lemma 2 we have

$\alpha_i^! \in F$ for $i = 1, \dots, n$, so $F \ni \bigwedge_{i=1}^n \alpha_i^!$
 $= \left(\bigvee_{i=1}^n \alpha_i \right)^! \neq \emptyset$ (because F is a proper filter)

Let $G = \{\alpha^! : \alpha \in A\}$. G has the FIP

because of that). So by lemma 1

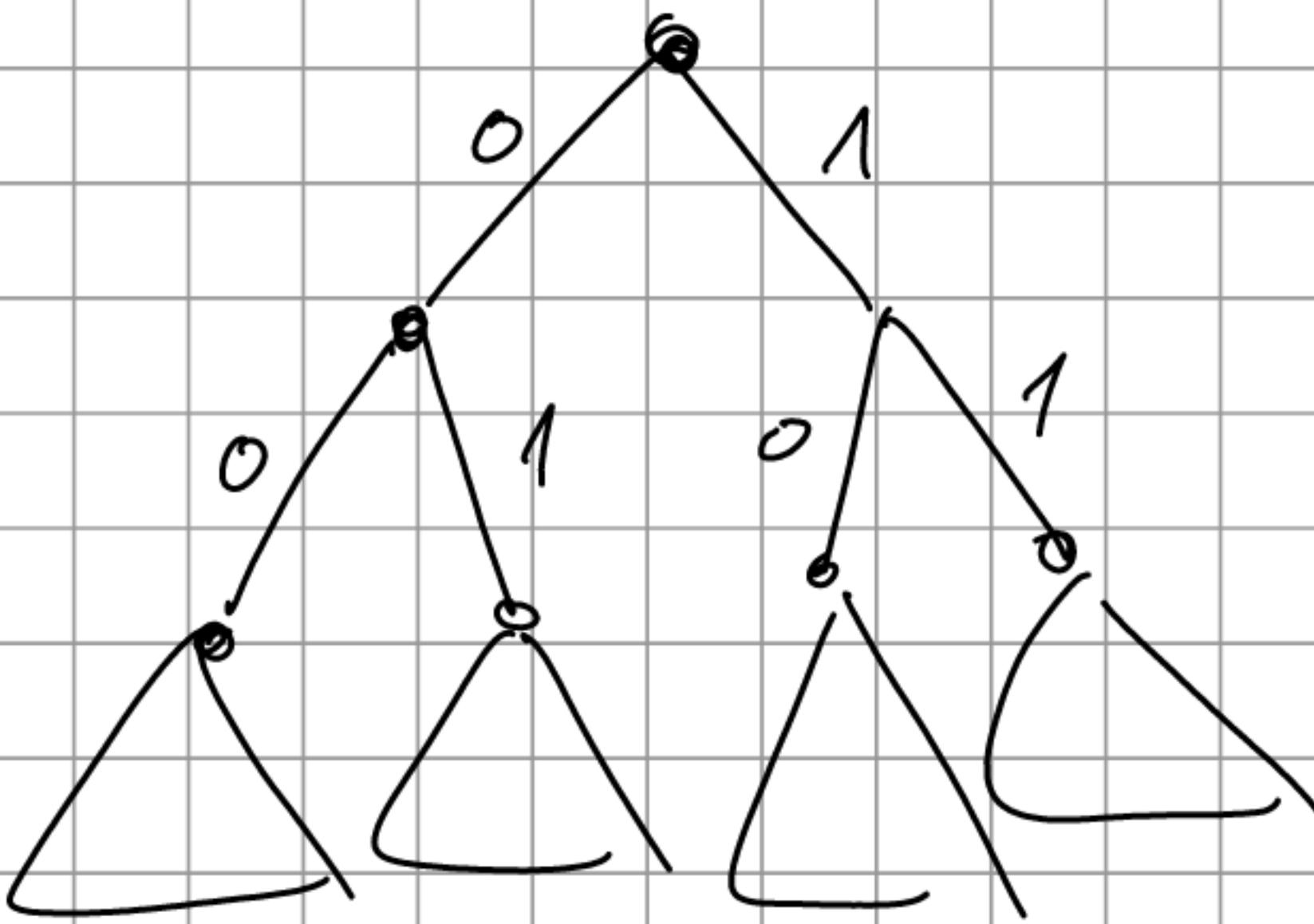
we can extend G

to a proper filter which can be extended to an ultrafilter (by ex. 5). Let's call

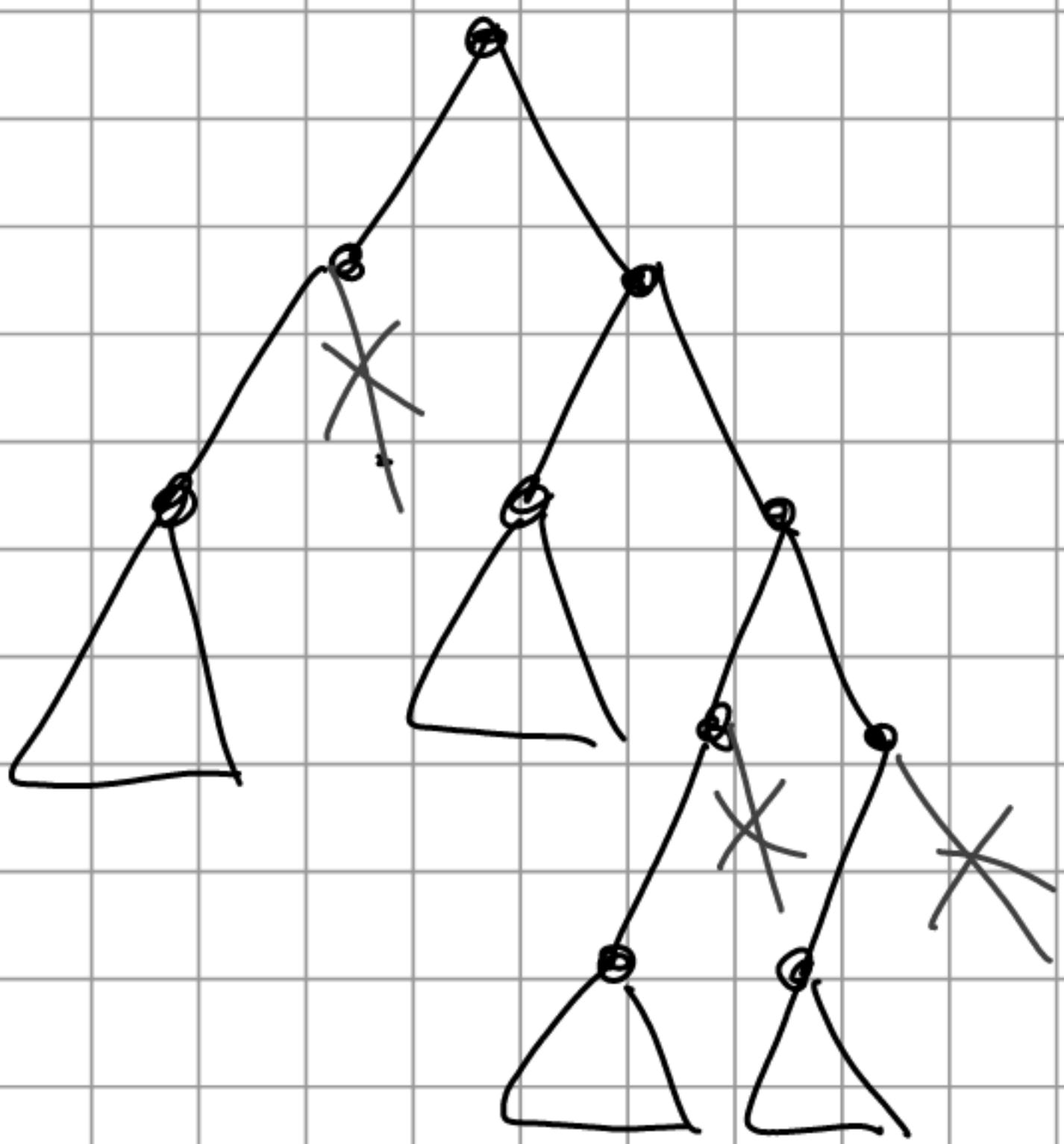
this ultrafilter \mathcal{U} . For every $a \in A$ $a' \in \mathcal{U}$,
so $a \notin \mathcal{U}$, so $\mathcal{U} \notin \bigcup_{a \in A} [a] = S(B)$.

Roof 2)

$B = \{b_1, b_2, \dots\}$. We want to construct
a homeomorphism $f: B \rightarrow \overline{\mathbb{D}}$, where
 $\mathbb{D} \subseteq \mathbb{C} = \{0, 1\}^{\mathbb{N}}$. We can think of
Cantor set as a binary tree:



A point in \mathbb{C} is an infinite branch
and a closed set is a tree
with some subtrees "cut":



Let's define f . For $F \in S(B)$ define

an auxiliary function $\sigma_F : B \rightarrow \{0, 1\}$

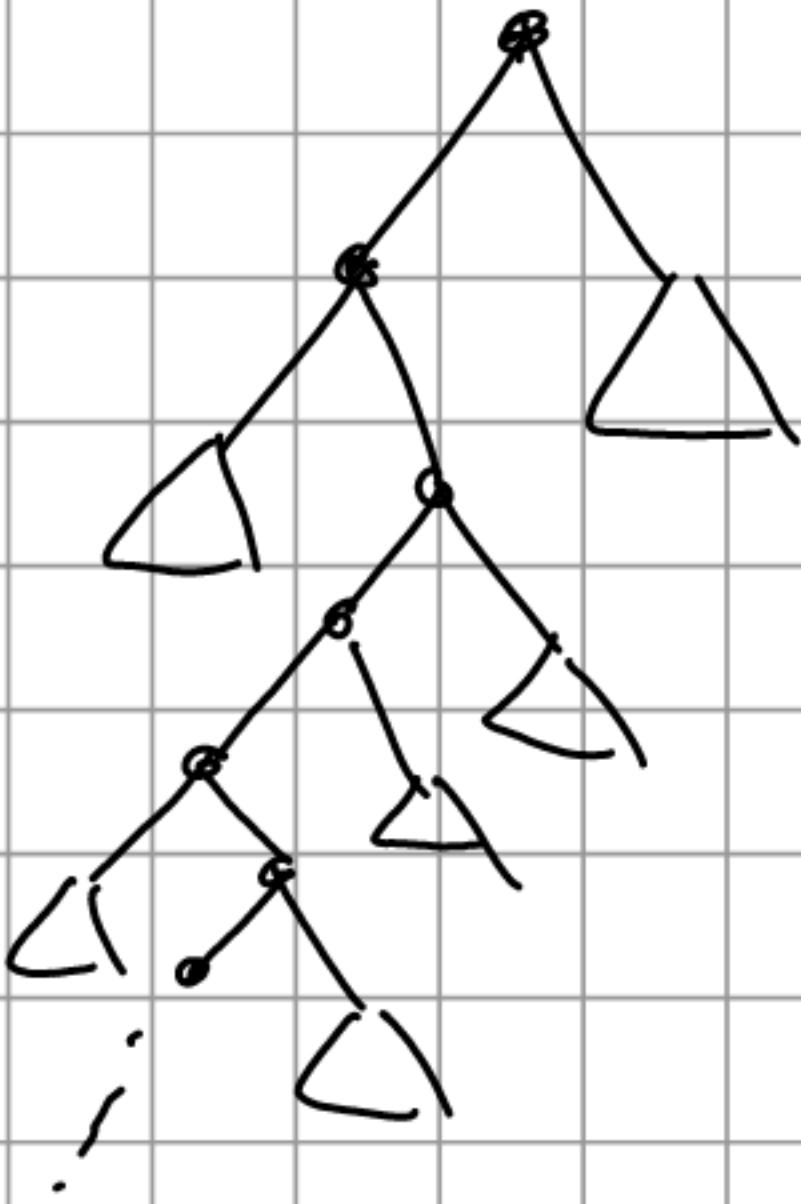
$$\text{with a formula } \sigma_F(b) = \begin{cases} 1 & b \in F \\ 0 & b' \in F \end{cases}$$

By lemma 2 this is a correct definition. Now:

$$f(F) = \langle \sigma_F(b_1), \sigma_F(b_2), \dots \rangle$$

First let's see that $\text{Rng}(f)$ is a closed subset of \mathbb{C} .

To see that let's take any $x \in \overline{\text{Rng}(f)}$. x is some branch



$$x = (x_1, x_2, \dots)$$

We want to show that $x \in \text{Rng}(f)$,
i.e. there is some $F \in S(B)$ s.t.

$$\sigma_F(b_i) = x_i \text{ for } i=1, 2, \dots \quad (\text{A.a.})$$

suppose there's no such F . What it

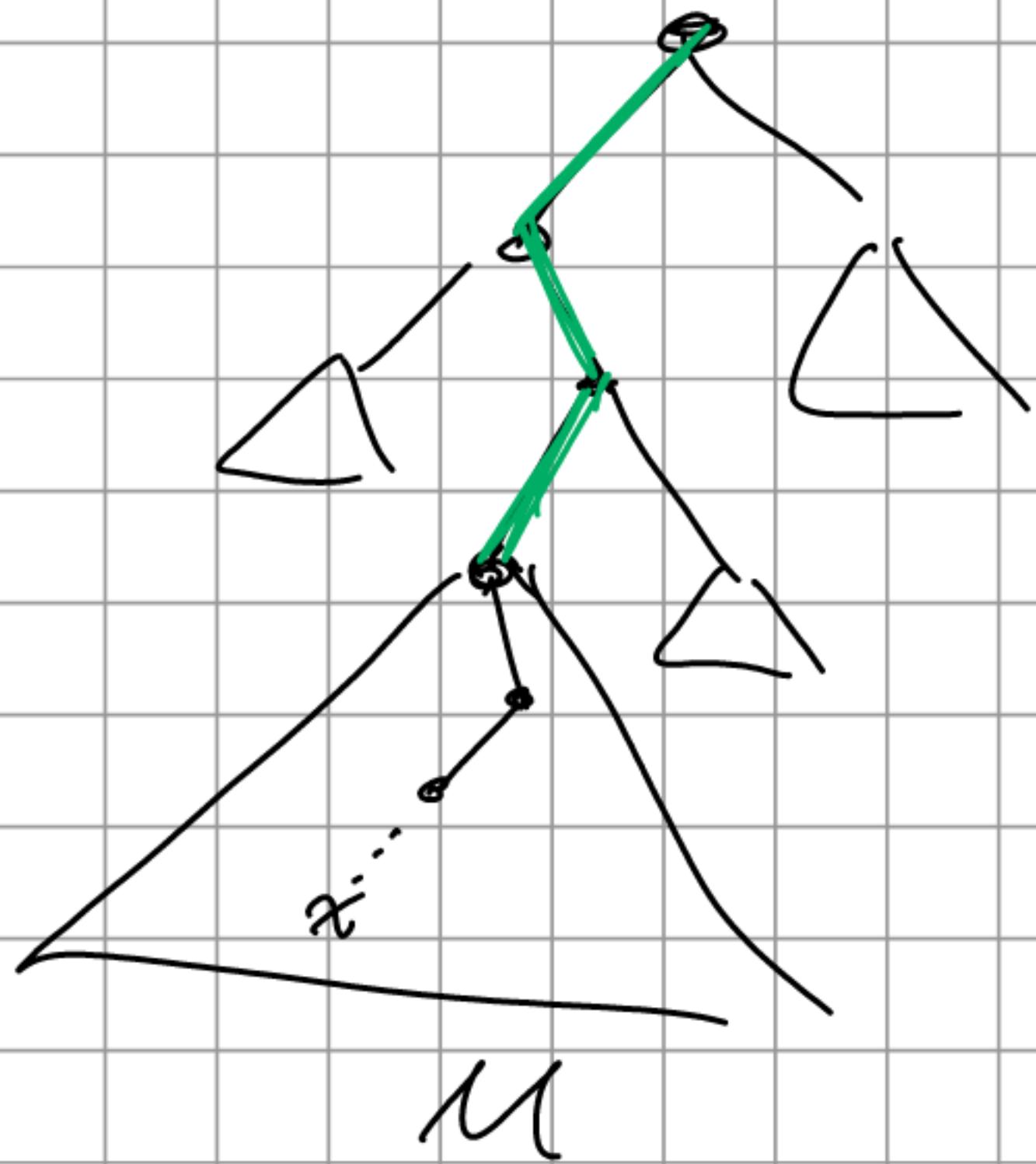
means is that for some prefix x_1, \dots, x_n

we have $\bigwedge_{i: x_i=1}^n b_i = \emptyset$ (if there's no

such prefix, then by lemma 1 there

is F s.t. $f(F) = x$). But "prefix"

is a subtree, which is a neighbour host
of x :



Moreover this neighbourhood has an

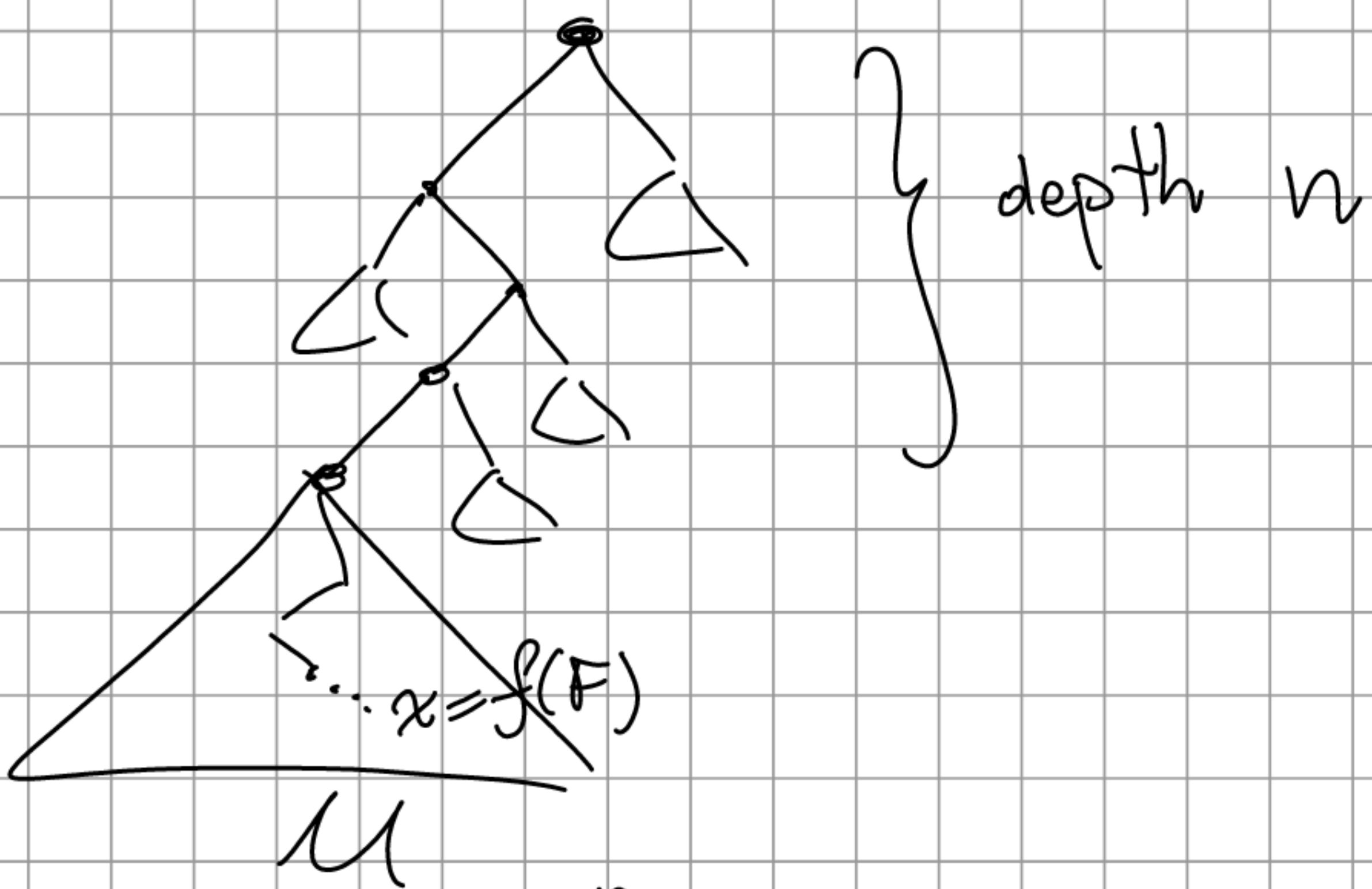
empty intersection with $\text{Rng}(f)$!

↙

Sidenote: I just realised that $S(B)$ is compact, so it's sufficient to show that f is a continuous 1-1 function, then we know that $f[S(B)] = D$ is compact and f is a homeomorphism. I'll leave the proof because it's nice in my opinion.

We'll show that f is continuous (1-1 is obvious).

Take $F \in S(B)$ and any base open neighbourhood \mathcal{U} of $f(F) = x$



Now let $b = \bigwedge_{i: b_i \in F}^n b_i$, $N = [b]$.
Then $f[N] \subseteq \mathcal{U}$.

This is trivial, because for every $F' \in N$ we have $\sigma_{F'}(b_i) = x_i$ where, $i = 1, \dots, n$, which means that $f(F')$ is in the subtree \mathcal{U} . ■

Lx.5 Every proper filter F in \mathcal{B} extends to an ultrafilter.

Proof let F be a proper filter in \mathcal{B} .

Let $\mathcal{F} = \{F' : F \subseteq F', F' \text{ is a prop. filter}\}$.

Take any chain $C \subseteq \mathcal{F}$. Then $\cup C$ is a proper filter. Take any $a, b \in C$. Then there

is some filter $F \in C$ s.t. $a, b \in F$,

so $a \wedge b \in F \Rightarrow a \wedge b \in C$. Of course

$\emptyset \notin C$ because \emptyset is not in any $F \in C$.

It is closed under supersets because every filter in C is.

By Zorn lemma we conclude that

there is a maximal $M \in \mathcal{F}$. It is

an ultrafilter. Suppose not: then there's

a prop.fil. $M' \supseteq M$, but also $F \subseteq M'$, so

$M' \in \mathcal{F}$ which contradicts maximality of M .