

1. INTRODUCTION

There will be something!

2. PRELIMINARIES

2.1. Descriptive set theory.

Definition 2.1. Suppose X is a topological space and $A \subseteq X$. We say that A is *meagre* in X if $A = \bigcup_{n \in \mathbb{N}} A_n$, where A_n are nowhere dense subsets of X (i.e. $\text{Int}(\bar{A}_n) = \emptyset$).

Definition 2.2. We say that A is *comeagre* in X if it is a complement of a meagre set. Equivalently, a set is comeagre if and only if it contains a countable intersection of open dense sets.

Every countable set is meagre in any T_1 space, so, for example, \mathbb{Q} is meagre in \mathbb{R} (although it is dense), which means that the set of irrationals is comeagre. Another example is...

Definition 2.3. We say that a topological space X is a *Baire space* if every comeagre subset of X is dense in X (equivalently, every meagre set has empty interior).

Definition 2.4. Suppose X is a Baire space. We say that a property P holds *generically* for a point $x \in X$ if $\{x \in X \mid P \text{ holds for } x\}$ is comeagre in X .

Definition 2.5. Let X be a nonempty topological space and let $A \subseteq X$. The *Banach-Mazur game* of A , denoted as $G^{**}(A)$ is defined as follows: Players *I* and *II* take turns in playing nonempty open sets $U_0, V_0, U_1, V_1, \dots$ such that $U_0 \supseteq V_0 \supseteq U_1 \supseteq V_1 \supseteq \dots$. We say that player *II* wins the game if $\bigcap_n V_n \subseteq A$.

There is an important theorem on the Banach-Mazur game: A is comeagre if and only if *II* can always choose sets V_0, V_1, \dots such that it wins. Before we prove it we need to define notions necessary to formalise and prove the theorem.

Definition 2.6. T is the *tree of all legal positions* in the Banach-Mazur game $G^{**}(A)$ when T consists of all finite sequences (W_0, W_1, \dots, W_n) , where W_i are nonempty open sets such that $W_0 \supseteq W_1 \supseteq \dots \supseteq W_n$. In other words, T is a pruned tree on $\{W \subseteq X \mid W \text{ is open nonempty}\}$.

Definition 2.7. We say that σ is a *pruned subtree* of the tree of all legal positions T if $\sigma \subseteq T$, for any $(W_0, W_1, \dots, W_n) \in \sigma, n \geq 0$ there is a W such that $(W_0, W_1, \dots, W_n, W) \in \sigma$ (it simply means that there's no finite branch in σ) and $(W_0, W_1, \dots, W_{n-1}) \in \sigma$ (every node on a branch is in σ).

Definition 2.8. Let σ be a pruned subtree of the tree of all legal positions T . By $[\sigma]$ we denote *the set of all infinite branches of σ* , i.e. infinite sequences (W_0, W_1, \dots) such that $(W_0, W_1, \dots, W_n) \in \sigma$ for any $n \in \mathbb{N}$.

Definition 2.9. A strategy for II in $G^{**}(A)$ is a pruned subtree $\sigma \subseteq T$ such that

- (i) σ is nonempty,
- (ii) if $(U_0, V_0, \dots, U_n, V_n) \in \sigma$, then for all open nonempty $U_{n+1} \subseteq V_n$, $(U_0, V_0, \dots, U_n, V_n, U_{n+1}) \in \sigma$,
- (iii) if $(U_0, V_0, \dots, U_n) \in \sigma$, then for a unique V_n , $(U_0, V_0, \dots, U_n, V_n) \in \sigma$.

Intuitively, a strategy σ works as follows: I starts playing U_0 as any open subset of X , then II plays unique (by (iii)) V_0 such that $(U_0, V_0) \in \sigma$. Then I responds by playing any $U_1 \subseteq V_0$ and II plays unique V_1 such that $(U_0, V_0, U_1, V_1) \in \sigma$, etc.

Definition 2.10. A strategy σ is a *winning strategy for II* if for any game $(U_0, V_0, \dots) \in [\sigma]$ player II wins, i.e. $\bigcap_n V_n \subseteq A$.

Now we can state the key theorem.

Theorem 2.11 (Banach-Mazur, Oxtoby). *Let X be a nonempty topological space and let $A \subseteq X$. Then A is comeagre $\Leftrightarrow II$ has a winning strategy in $G^{**}(A)$.*

In order to prove it we add an auxiliary definition and lemma.

Definition 2.12. Let $S \subseteq \sigma$ be a pruned subtree of tree of all legal positions T and let $p = (U_0, V_0, \dots, V_n) \in S$. We say that S is *comprehensive for p* if the family $\mathcal{V}_p = \{V_{n+1} \mid (U_0, V_0, \dots, V_n, U_{n+1}, V_{n+1}) \in S\}$ (it may be that $n = -1$, which means $p = \emptyset$) is pairwise disjoint and $\bigcup \mathcal{V}_p$ is dense in V_n (where we think that $V_{-1} = X$). We say that S is *comprehensive* if it is comprehensive for each $p = (U_0, V_0, \dots, V_n) \in S$.

Fact 2.13. *If σ is a winning strategy for II then there exists a nonempty comprehensive $S \subseteq \sigma$.*

Proof. We construct S recursively as follows:

- (1) $\emptyset \in S$,
- (2) if $(U_0, V_0, \dots, U_n) \in S$, then $(U_0, V_0, \dots, U_n, V_n) \in S$ for the unique V_n given by the strategy σ ,
- (3) let $p = (U_0, V_0, \dots, V_n) \in S$. For a possible player I 's move $U_{n+1} \subseteq V_n$ let U_{n+1}^* be the unique set player II would respond with by σ . Now, by Zorn's Lemma, let \mathcal{U}_p be a maximal collection of nonempty open subsets $U_{n+1} \subseteq V_n$ such that the set $\{U_{n+1}^* \mid U_{n+1} \in \mathcal{U}_p\}$ is pairwise disjoint. Then put in S all $(U_0, V_0, \dots, V_n, U_{n+1})$ such that $U_{n+1} \in \mathcal{U}_p$. This way S is comprehensive for p : the family $\mathcal{V}_p = \{V_{n+1} \mid (U_0, V_0, \dots, V_n, U_{n+1}, V_{n+1}) \in S\}$ is exactly $\{U_{n+1}^* \mid U_{n+1} \in \mathcal{U}_p\}$, which is pairwise disjoint and $\bigcup \mathcal{V}_p$ is obviously dense in V_n by the maximality of \mathcal{U}_p – if there was any open set $\tilde{U}_{n+1} \subseteq V_n$ disjoint from $\bigcup \mathcal{V}_p$, then $\tilde{U}_{n+1}^* \subseteq \tilde{U}_{n+1}$ would be also disjoint from

$\bigcup \mathcal{V}_p$, so the family $\mathcal{U}_p \cup \{\tilde{U}_{n+1}\}$ would violate the maximality of \mathcal{U}_p . \square

Lemma 2.14. *Let S be a nonempty comprehensive pruned subtree of a strategy σ . Then:*

- (i) *For any open $V_n \subseteq X$ there is at most one $p = (U_0, V_0, \dots, U_n, V_n) \in S$.*
- (ii) *Let $S_n = \{V_n \mid (U_0, V_0, \dots, V_n) \in S\}$ for $n \in \mathbb{N}$ (i.e. S_n is a family of all possible choices player II can make in its n -th move according to S). Then $\bigcup S_n$ is open and dense in X .*
- (iii) *S_n is a family of pairwise disjoint sets.*

Proof. (i): Suppose that there are some $p = (U_0, V_0, \dots, U_n, V_n)$, $p' = (U'_0, V'_0, \dots, U'_n, V'_n)$ such that $V_n = V'_n$ and $p \neq p'$. Let k be the smallest index such that those sequences differ. We have two possibilities:

- $U_k = U'_k$ and $V_k \neq V'_k$ – this cannot be true simply by the fact that S is a subset of a strategy (so V_k is unique for U_k).
- $U_k \neq U'_k$: by the comprehensiveness of S we know that for $q = (U_0, V_0, \dots, U_{k-1}, V_{k-1})$ the set \mathcal{V}_q is pairwise disjoint. Thus $V_k \cap V'_k = \emptyset$, because $V_k, V'_k \in \mathcal{V}_q$. But this leads to a contradiction – V_n cannot be a nonempty subset of both V_k, V'_k .

(ii): The lemma is proved by induction on n . For $n = 0$ it follows trivially from the definition of comprehensiveness. Now suppose the lemma is true for n . Then the set $\bigcup_{V_n \in S_n} \bigcup \mathcal{V}_{p_{V_n}}$ (where p_{V_n} is given uniquely from (i)) is dense and open in X by the induction hypothesis. But $\bigcup S_{n+1}$ is exactly this set, thus it is dense and open in X .

(iii): We will prove it by induction on n . Once again, the case $n = 0$ follows from the comprehensiveness of S . Now suppose that the sets in S_n are pairwise disjoint. Take some $x \in V_{n+1} \in S_{n+1}$. Of course $\bigcup S_n \supseteq \bigcup S_{n+1}$, thus by the inductive hypothesis $x \in V_n$ for the unique $V_n \in S_n$. It must be that $V_{n+1} \in \mathcal{V}_{p_{V_n}}$, because V_n is the only superset of V_{n+1} in S_n . But $\mathcal{V}_{p_{V_n}}$ is disjoint, so there is no other $V'_{n+1} \in \mathcal{V}_{p_{V_n}}$ such that $x \in V'_{n+1}$. Moreover, there is no such set in $S_{n+1} \setminus \mathcal{V}_{p_{V_n}}$, because those sets are disjoint from V_n . Hence there is no $V'_{n+1} \in S_{n+1}$ other than V_n such that $x \in V'_{n+1}$. We've chosen x and V_{n+1} arbitrarily, so S_{n+1} is pairwise disjoint. \square

Now we can move to the proof of the Banach-Mazur theorem.

Proof of theorem 2.11. \Rightarrow : Let (A_n) be a sequence of dense open sets with $\bigcap_n A_n \subseteq A$. The simply II plays $V_n = U_n \cap A_n$, which is nonempty by the denseness of A_n .

\Leftarrow : Suppose II has a winning strategy σ . We will show that A is comeagre. Take a comprehensive $S \subseteq \sigma$. We claim that $\mathcal{S} = \bigcap_n \bigcup S_n \subseteq A$. By the lemma 2.14, (ii) sets $\bigcup S_n$ are open and dense, thus A must be comeagre. Now we prove the claim towards contradiction.

Suppose there is $x \in \mathcal{S} \setminus A$. By the lemma 2.14, (iii) for any n there is unique $x \in V_n \in \mathcal{S}_n$. It follows that $p_{V_0} \subset p_{V_1} \subset \dots$. Now the game $(U_0, V_0, U_1, V_1, \dots) = \bigcup_n p_{V_n} \in [S] \subseteq [\sigma]$ is not winning for player II, which contradicts the assumption that σ is a winning strategy. \square

Corollary 2.15. *If we add a constraint to the Banach-Mazur game such that players can only choose basic open sets, then the theorem 2.11 still suffices.*

Proof. If one adds the word *basic* before each occurrence of word *open* in previous proofs and theorems then they all will still be valid (except for \Rightarrow , but its an easy fix – take V_n a basic open subset of $U_n \cap A_n$). \square

This corollary will be important in using the theorem in practice – it’s much easier to work with basic open sets rather than any open sets.

2.2. Category theory. In this section we will give a short introduction to the notions of category theory that will be necessary to generalize the key result of the paper.

We will use a standard notation. If the reader is interested in detailed introduction to the category theory, then it’s recommended to take a look at [2]. Here we will shortly describe the standard notation.

A *category* \mathcal{C} consists of the collection of objects (denoted as $\text{Obj}(\mathcal{C})$, but most often simply as \mathcal{C}) and collection of *morphisms* $\text{Mor}(A, B)$ between each pair of objects $A, B \in \mathcal{C}$. We require that for each morphisms $f : B \rightarrow C$, $g : A \rightarrow B$ there is a morphism $f \circ g : A \rightarrow C$. For every $A \in \mathcal{C}$ there is an *identity morphism* id_A such that for any morphism $f \in \text{Mor}(A, B)$ it follows that $f \circ \text{id}_A = \text{id}_B \circ f$.

We say that $f : A \rightarrow B$ is *isomorphism* if there is (necessarily unique) morphism $g : B \rightarrow A$ such that $g \circ f = \text{id}_A$ and $f \circ g = \text{id}_B$. Automorphism is an isomorphism where $A = B$.

A *functor* is a “homeomorphism“ of categories. We say that $F : \mathcal{C} \rightarrow \mathcal{D}$ is a functor from category \mathcal{C} to category \mathcal{D} if it associates each object $A \in \mathcal{C}$ with an object $F(A) \in \mathcal{D}$, associates each morphism $f : A \rightarrow B$ in \mathcal{C} with a morphism $F(f) : F(A) \rightarrow F(B)$. We also require that $F(\text{id}_A) = \text{id}_{F(A)}$ and that for any (compatible) morphisms f, g in \mathcal{C} $F(f \circ g) = F(f) \circ F(g)$.

In category theory we distinguish *covariant* and *contravariant* functors. Here, we only consider *covariant functors*, so we will simply say *functor*.

Fact 2.16. *Functor $F : \mathcal{C} \rightarrow \mathcal{D}$ maps isomorphism $f : A \rightarrow B$ in \mathcal{C} to the isomorphism $F(f) : F(A) \rightarrow F(B)$ in \mathcal{D} .*

Notion that will be very important for us is a “morphism of functors“ which is called *natural transformation*.

Definition 2.17. Let F, G be functors between the categories \mathcal{C}, \mathcal{D} . A *natural transformation* τ is function that assigns to each object A of \mathcal{C} a

morphism τ_A in $\text{Mor}(F(A), G(A))$ such that for every morphism $f : A \rightarrow B$ in \mathcal{C} the following diagram commutes:

$$\begin{array}{ccccc} A & & F(A) & \xrightarrow{\tau_A} & G(A) \\ \downarrow f & & \downarrow F(f) & & \downarrow G(f) \\ B & & F(B) & \xrightarrow{\tau_B} & G(B) \end{array}$$

Definition 2.18. In category theory, a *diagram* of type \mathcal{J} in category \mathcal{C} is a functor $D: \mathcal{J} \rightarrow \mathcal{C}$. \mathcal{J} is called the *index category* of D . In other words, D is of *shape* \mathcal{J} .

For example, $\mathcal{J} = \{-1 \leftarrow 0 \rightarrow 1\}$, then a diagram $D: \mathcal{J} \rightarrow \mathcal{C}$ is called a *cospan*. For example, if A, B, C are objects of \mathcal{C} and $f \in \text{Mor}(C, A)$, $g \in \text{Mor}(C, B)$, then the following diagram is a cospan:

$$\begin{array}{ccc} A & & B \\ & \swarrow f & \nearrow g \\ & C & \end{array}$$

From now we omit explicit definition of the index category, as it is easily referable from a picture.

Definition 2.19. Let A, B, C, D be objects in the category \mathcal{C} with morphisms $e: C \rightarrow A$, $f: C \rightarrow B$, $g: A \rightarrow D$, $h: B \rightarrow D$ such that $g \circ e = h \circ f$. Then the following diagram:

$$\begin{array}{ccc} & D & \\ g \nearrow & & \nwarrow h \\ A & & B \\ f \nwarrow & & \nearrow e \\ & C & \end{array}$$

is called a *pushout* diagram

Definition 2.20. The *cospan category* of category \mathcal{C} , referred to as $\text{Cospan}(\mathcal{C})$, is the category of cospan diagrams of \mathcal{C} , where morphisms between two cospans are normal transformations of the underlying functors.

We define *pushout category* analogously and call it $\text{Pushout}(\mathcal{C})$.

TODO: dodać tu przykład?

3. FRAÏSSÉ CLASSES

In this section we will take a closer look at classes of finitely generated structures with some characteristic properties. More specifically, we will describe a concept developed by a French mathematician Roland Fraïssé called Fraïssé limit.

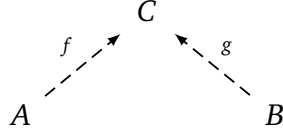
3.1. Definitions.

Definition 3.1. Let L be a signature and M be an L -structure. The *age* of M is the class \mathcal{K} of all finitely generated structures that embeds into M . The age of M is also associated with class of all structures embeddable in M up to isomorphism.

Definition 3.2. We say that M has *countable age* when its age has countably many isomorphism types of finitely generated structures.

Definition 3.3. Let \mathcal{K} be a class of finitely generated structures. \mathcal{K} has *hereditary property (HP)* if for any $A \in \mathcal{K}$, any finitely generated substructure B of A it holds that $B \in \mathcal{K}$.

Definition 3.4. Let \mathcal{K} be a class of finitely generated structures. We say that \mathcal{K} has *joint embedding property (JEP)* if for any $A, B \in \mathcal{K}$ there is a structure $C \in \mathcal{K}$ such that both A and B embed in C .



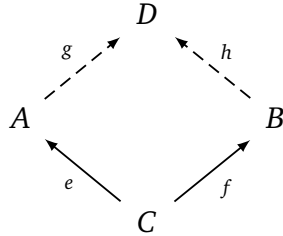
In terms of category theory, this is a *span* in category \mathcal{K} .

Fraïssé has shown fundamental theories regarding age of a structure, one of them being the following one:

Fact 3.5. Suppose L is a signature and \mathcal{K} is a nonempty finite or countable set of finitely generated L -structures. Then \mathcal{K} has the HP and JEP if and only if \mathcal{K} is the age of some finite or countable structure.

Beside the HP and JEP Fraïssé has distinguished one more property of the class \mathcal{K} , namely amalgamation property.

Definition 3.6. Let \mathcal{K} be a class of finitely generated L -structures. We say that \mathcal{K} has the *amalgamation property (AP)* if for any $A, B, C \in \mathcal{K}$ and embeddings $e: C \rightarrow A, f: C \rightarrow B$ there exists $D \in \mathcal{K}$ together with embeddings $g: A \rightarrow D$ and $h: B \rightarrow D$ such that $g \circ e = h \circ f$.



In terms of category theory, amalgamation over some structure C is a pushout diagram.

Definition 3.7. Let M be an L -structure. M is *ultrahomogeneous* if every isomorphism between finitely generated substructures of M extends to an automorphism of M .

Having those definitions we can provide the main Fraïssé theorem.

Theorem 3.8 (Fraïssé theorem). *Let L be a countable language and let \mathcal{K} be a nonempty countable set of finitely generated L -structures which has HP, JEP and AP. Then \mathcal{K} is the age of a countable, ultrahomogeneous L -structure M . Moreover, M is unique up to isomorphism. We say that M is a Fraïssé limit of \mathcal{K} and denote this by $M = \text{Flim}(\mathcal{K})$.*

This is a well known theorem. One can read a proof of this theorem in Wilfrid Hodges' classical book *Model Theory* [1]. In the proof of this theorem appears another, equally important 3.10.

Definition 3.9. We say that an L -structure M is *weakly ultrahomogeneous* if for any A, B finitely generated substructures of M such that $A \subseteq B$ and an embedding $f: A \rightarrow M$ there is an embedding $g: B \rightarrow M$ which extends f .

$$\begin{array}{ccc} A & \xrightarrow{f} & D \\ \subseteq \downarrow & \nearrow g & \\ B & & \end{array}$$

Lemma 3.10. *A countable structure is ultrahomogeneous if and only if it is weakly ultrahomogeneous.*

This lemma will play a major role in the later parts of the paper. Weak ultrahomogeneity is an easier and more intuitive property and it will prove useful when recursively constructing the generic automorphism of a Fraïssé limit.

3.2. Random graph. In this section we'll take a closer look on a class of finite unordered graphs, which is a Fraïssé class.

The language of unordered graphs L consists of a single binary relational symbol E . If G is an L -structure then we call it a *graph*, and its elements *vertices*. If for some vertices $u, v \in G$ we have $G \models uEv$ then we say that there is an *edge* connecting u and v . If $G \models \forall x \forall y (xEy \leftrightarrow yEx)$ then we say that G is an *unordered graph*. From now on we omit the word *unordered* and graphs as *unordered*.

Proposition 3.11. *Let \mathcal{G} be the class of all finite graphs. \mathcal{G} is a Fraïssé class.*

Proof. \mathcal{G} is of course countable (up to isomorphism) and has the HP (graph substructure is also a graph). It has JEP: having two finite graphs G_1, G_2 take their disjoint union $G_1 \sqcup G_2$ as the extension of them both. \mathcal{G} has the AP. Having graphs A, B, C , where B and C are supergraphs of A , we can assume without loss of generality, that $(B \setminus A) \cap (C \setminus A) = \emptyset$. Then $A \sqcup (B \setminus A) \sqcup (C \setminus A)$ is the graph we're looking for (with edges as in B and C and without any edges between $B \setminus A$ and $C \setminus A$). \square

Definition 3.12. The *random graph* is the Fraïssé limit of the class of finite graphs \mathcal{G} denoted by $\Gamma = \text{Flim}(\mathcal{G})$.

The concept of the random graph emerges independently in many fields of mathematics. For example, one can construct the graph by choosing at random for each pair of vertices if they should be connected or not. It turns out that the graph constructed this way is isomorphic to the random graph with probability 1.

The random graph Γ has one particular property that is unique to the random graph.

Fact 3.13 (random graph property). *For each finite disjoint $X, Y \subseteq \Gamma$ there exists $v \in \Gamma \setminus (X \cup Y)$ such that $\forall u \in X (vEu)$ and $\forall u \in Y (\neg vEu)$.*

Proof. Take any finite disjoint $X, Y \subseteq \Gamma$. Let G_{XY} be the subgraph of Γ induced by the $X \cup Y$. Let $H = G_{XY} \cup \{w\}$, where w is a new vertex that does not appear in G_{XY} . Also, w is connected to all vertices of G_{XY} that come from X and to none of those that come from Y . This graph is of course finite, so it is embeddable in Γ . Without loss of generality assume that this embedding is simply inclusion. Let f be the partial isomorphism from $X \sqcup Y$ to H , with X and Y projected to the part of H that come from X and Y respectively. By the ultrahomogeneity of Γ this isomorphism extends to an automorphism $\sigma \in \text{Aut}(\Gamma)$. Then $v = \sigma^{-1}(w)$ is the vertex we sought. \square

Fact 3.14. *If a countable graph G has the random graph property, then it is isomorphic to the random graph Γ .*

Proof. Enumerate vertices of both graphs: $\Gamma = \{a_1, a_2, \dots\}$ and $G = \{b_1, b_2, \dots\}$. We will construct a chain of partial isomorphisms $f_n: \Gamma \rightarrow G$ such that $\emptyset = f_0 \subseteq f_1 \subseteq f_2 \subseteq \dots$ and $a_n \in \text{dom}(f_n)$ and $b_n \in \text{rng}(f_n)$.

Suppose we have f_n . We seek $b \in G$ such that $f_n \cup \{\langle a_{n+1}, b \rangle\}$ is a partial isomorphism. If $a_{n+1} \in \text{dom } f_n$, then simply $b = f_n(a_{n+1})$. Otherwise, let $X = \{a \in \Gamma \mid aE_\Gamma a_{n+1}\} \cap \text{dom } f_n, Y = X^c \cap \text{dom } f_n$, i.e. X are vertices of $\text{dom } f_n$ that are connected with a_{n+1} in Γ and Y are those vertices that are not connected with a_{n+1} . Let b be a vertex of G that is connected to all vertices of $f_n[X]$ and to none $f_n[Y]$ (it exists by the random graph property). Then $f_n \cup \{\langle a_{n+1}, b \rangle\}$ is a partial isomorphism. We find a for the b_{n+1} in the similar manner, so that $f_{n+1} = f_n \cup \{\langle a_{n+1}, b \rangle, \langle a, b_{n+1} \rangle\}$ is a partial isomorphism.

Finally, $f = \bigcup_{n=0}^{\infty} f_n$ is an isomorphism between Γ and G . Take any $a, b \in \Gamma$. Then for some big enough n we have that $aE_\Gamma b \Leftrightarrow f_n(a)E_G f_n(b) \Leftrightarrow f(a)E_G f(b)$. \square

Using this fact one can show that the graph constructed in the probabilistic manner is in fact isomorphic to the random graph Γ .

Definition 3.15. We say that a Fraïssé class \mathcal{K} has *weak Hrushovski property (WHP)* if for every $A \in \mathcal{K}$ and an isomorphism of its substructures $p: A \rightarrow A$ (also called a partial automorphism of A), there is some

$B \in \mathcal{K}$ such that p can be extended to an automorphism of B , i.e. there is an embedding $i: A \rightarrow B$ and a $\bar{p} \in \text{Aut}(B)$ such that the following diagram commutes:

$$\begin{array}{ccc} B & \xrightarrow{\bar{p}} & B \\ \uparrow i & & \uparrow i \\ A & \xrightarrow{p} & A \end{array}$$

Proposition 3.16. *The class of finite graphs \mathcal{G} has the weak Hrushovski property.*

Proof. It may be there some day, but it may not! □

3.3. Canonical amalgamation.

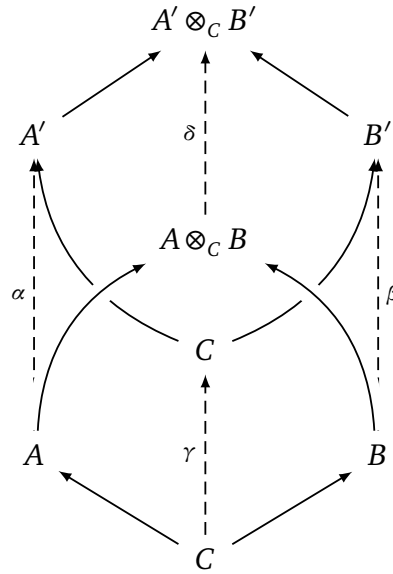
Definition 3.17. Let \mathcal{K} be a class finitely generated L -structures. We say that \mathcal{K} has *canonical amalgamation* if for every $C \in \mathcal{K}$ there is a functor $\otimes_C: \text{Cospan}(\mathcal{K}) \rightarrow \text{Pushout}(\mathcal{K})$ such that it has the following properties:

- Let $A \leftarrow C \rightarrow B$ be a cospan. Then \otimes_C sends it to a pushout that preserves “the bottom“ structures and embeddings, i.e.:

$$\begin{array}{ccccc} & & & & A \otimes_C B \\ & & & & \swarrow \quad \searrow \\ & & & & \\ A & & B & \xrightarrow{A \otimes_C B} & A & & B \\ & \swarrow & \searrow & & \swarrow & \searrow \\ & C & & & C & & \\ & \swarrow & \searrow & & \swarrow & \searrow \\ & & & & & & \end{array}$$

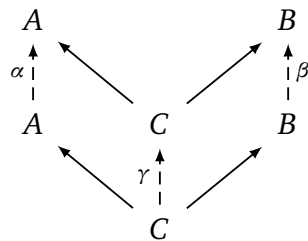
We have deliberately omitted names for embeddings of C . Of course, the functor has to take them into account, but this abuse of notation is convenient and should not lead into confusion.

- Let $A \leftarrow C \rightarrow B, A' \leftarrow C \rightarrow B'$ be cospans with a natural transformation given by $\alpha: A \rightarrow A', \beta: B \rightarrow B', \gamma: C \rightarrow C$. Then \otimes_C preserves the morphisms of when sending the natural transformation of those cospans to natural transformation of pushouts by adding the $\delta: A \otimes_C B \rightarrow A' \otimes_C B'$ morphism:

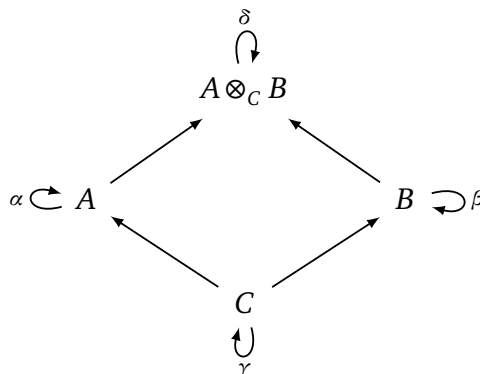


Theorem 3.18. *Let \mathcal{K} be a Fraïssé class of L -structures with canonical amalgamation. Then the class \mathcal{H} of L -structures with automorphism is a Fraïssé class.*

Proof. \mathcal{H} is obviously countable and has HP. It suffices to show that it has AP (JEP follows by taking C to be the empty structure). Take any $(A, \alpha), (B, \beta), (C, \gamma) \in \mathcal{H}$ such that (C, γ) embeds into (A, α) and (B, β) . Then α, β, γ yield an automorphism η (as a natural transformation) of a cospan:



Then, by the fact 2.16, $\otimes_C(\eta)$ is an automorphism of the pushout diagram:



TODO: ten diagram nie jest do końca taki jak trzeba, trzeba w zasadzie skopiować ten z definicji kanonicznej amalgamacji. Czy to nie będzie wyglądać źle?

This means that the morphism $\delta: A \otimes_C B \rightarrow A \otimes_C B$ has to be automorphism. Thus, by the fact that the diagram commutes, we have the amalgamation of (A, α) and (B, β) over (C, γ) in \mathcal{H} . \square

3.4. Graphs with automorphism. The language and theory of unordered graphs is fairly simple. We extend the language by one unary symbol σ and interpret it as an arbitrary automorphism on the graph structure. It turns out that the class of such structures is a Fraïssé class.

Proposition 3.19. *Let \mathcal{H} be the class of all finite graphs with an automorphism, i.e. structures in the language (E, σ) such that E is a symmetric relation and σ is an automorphism on the structure. \mathcal{H} is a Fraïssé class.*

Proof. Countability and HP are obvious, JEP follows by the same argument as in plain graphs. We need to show that the class has the amalgamation property.

Take any $(A, \alpha), (B, \beta), (C, \gamma) \in \mathcal{H}$ such that A embeds into B and C . Without the loss of generality we may assume that $B \cap C = A$ and $\alpha \subseteq \beta, \gamma$. Let D be the amalgamation of B and C over A as in the proof for the plain graphs. We will define the automorphism $\delta \in \text{Aut}(D)$ so it extends β and γ . We let $\delta \upharpoonright_B = \beta, \delta \upharpoonright_C = \gamma$. Let's check the definition is correct. We have to show that $(uE_D v \leftrightarrow \delta(u)E_D \delta(v))$ holds for any $u, v \in D$. We have two cases:

- $u, v \in X$, where X is either B or C . This case is trivial.
- $u \in B \setminus A, v \in C \setminus A$. Then $\delta(u) = \beta(u) \in B \setminus A$, similarly $\delta(v) = \gamma(v) \in C \setminus A$. This follows from the fact, that $\beta \upharpoonright_A = \alpha$, so for any $w \in A$ $\beta^{-1}(w) = \alpha^{-1}(w) \in A$, similarly for γ . Thus, from the construction of D , $\neg uE_D v$ and $\neg \delta(u)E_D \delta(v)$.

\square

The proposition says that there is a Fraïssé for the class \mathcal{H} of finite graphs with automorphisms. We shall call it (Π, σ) . Not surprisingly, Π is in fact a random graph.

Proposition 3.20. *The Fraïssé limit of \mathcal{H} interpreted as a plain graph (i.e. as a reduct to the language of graphs) is isomorphic to the random graph Γ .*

Proof. It is enough to show that $\Pi = \text{Flim}(\mathcal{H})$ has the random graph property. Take any finite disjoint $X, Y \subseteq \Pi$. Without the loss of generality assume that $X \cup Y$ is σ -invariant, i.e. $\sigma(v) \in X \cup Y$ for $v \in X \cup Y$. This assumption can be made because there are no infinite orbits in σ , which in turn is due to the fact that \mathcal{H} is the age of Π .

Let G_{XY} be the graph induced by $X \cup Y$. Take $H = G_{XY} \sqcup v$ as a supergraph of G_{XY} with one new vertex v connected to all vertices of X and to none of Y . By the proposition 3.16 we can extend H together with its partial isomorphism $\sigma \upharpoonright_{X \cup Y}$ to a graph R with automorphism τ . Once again, without the loss of generality we can assume that $R \subseteq \Pi$, because \mathcal{H} is the age of Π . But $R \upharpoonright_{G_{XY}}$ together with $\tau \upharpoonright_{G_{XY}}$ are isomorphic to the G_{XY} with $\sigma \upharpoonright_{G_{XY}}$.

Thus, by ultrahomogeneity of Π this isomorphism extends to an automorphism θ of (Π, σ) . Then $\theta(v)$ is the vertex that is connected to all vertices of X and none of Y , because $\theta[R \upharpoonright_X] = X$, $\theta[R \upharpoonright_Y] = Y$. \square

Theorem 3.21. *Let \mathcal{C} be a Fraïssé class of finite structures in a relational language L of some theory T . Let \mathcal{D} be a class of finite structures of the theory T in a relational language L with additional unary function symbol interpreted as an automorphism of the structure. If \mathcal{C} has the weak Hrushovski property and \mathcal{D} is a Fraïssé class then the Fraïssé limit of \mathcal{C} is isomorphic to the Fraïssé limit of \mathcal{D} reduced to the language L .*

Proof. Let $\Gamma = \text{Flim}(\mathcal{C})$ and $(\Pi, \sigma) = \text{Flim}(\mathcal{D})$. By the Fraïssé theorem 3.8 it suffices to show that the age of Π is \mathcal{C} and that it has the weak ultrahomogeneity in the class \mathcal{C} . The former comes easily, as for every structure $A \in \mathcal{C}$ we have the structure $(A, \text{id}_A) \in \mathcal{D}$, which means that the structure A embeds into Π . Also, if a structure $(B, \beta) \in \mathcal{D}$ embeds into \mathcal{D} , then $B \in \mathcal{C}$. Hence, \mathcal{C} is indeed the age of Π .

Now, take any structure $A, B \in \mathcal{C}$ such that $A \subseteq B$. Without the loss of generality assume that $A = B \cap \Pi$. Let \bar{A} be the smallest structure closed on the automorphism σ and containing A . It is finite, as \mathcal{C} is the age of Π . By the weak Hrushovski property, of \mathcal{C} let (\bar{B}, β) be a structure extending $(B \cup \bar{A}, \sigma \upharpoonright_{\bar{A}})$. Again, we may assume that $B \cup \bar{A} \subseteq \bar{B}$. Then, by the fact that Π is a Fraïssé limit of \mathcal{D} there is an embedding $f : (\bar{B}, \beta) \rightarrow (\Pi, \sigma)$ such that the following diagram commutes:

$$\begin{array}{ccccc} (A, \emptyset) & \xrightarrow{\subseteq} & (\bar{A}, \sigma \upharpoonright_{\bar{A}}) & \xrightarrow{\subseteq} & (\Pi, \sigma) \\ \subseteq \downarrow & & \downarrow \subseteq & \nearrow f & \\ (B, \sigma \upharpoonright_B) & \dashrightarrow^{\subseteq} & (\bar{B}, \beta) & & \end{array}$$

Then we simply get the following diagram:

$$\begin{array}{ccc}
A & \xrightarrow{\subseteq} & \Pi \\
\subseteq \downarrow & \nearrow f \upharpoonright_B & \\
B & &
\end{array}$$

which proves that Π is indeed a weakly ultrahomogeneous structure in \mathcal{C} . Hence, it is isomorphic to Γ . \square

4. CONJUGACY CLASSES IN AUTOMORPHISM GROUPS

Let M be a countable L -structure. We define a topology on the $G = \text{Aut}(M)$: for any finite function $f : M \rightarrow M$ we have a basic open set $[f]_G = \{g \in G \mid f \subseteq g\}$.

4.1. Prototype: pure set. In this section, $M = (M, =)$ is an infinite countable set (with no structure beyond equality).

Proposition 4.1. *If $f_1, f_2 \in \text{Aut}(M)$, then f_1 and f_2 are conjugate if and only if for each $n \in \mathbb{N} \cup \{\aleph_0\}$, f_1 and f_2 have the same number of orbits of size n .*

Theorem 4.2. *Let $\sigma \in \text{Aut}(M)$ be an automorphism with no infinite orbit and with infinitely many orbits of size n for every $n > 0$. Then the conjugacy class of σ is comeagre in $\text{Aut}(M)$.*

Proof. We will show that the conjugacy class of σ is an intersection of countably many comeagre sets.

Let $A_n = \{\alpha \in \text{Aut}(M) \mid \alpha \text{ has infinitely many orbits of size } n\}$. This set is comeagre for every $n > 0$. Indeed, we can represent this set as an intersection of countable family of open dense sets. Let $B_{n,k}$ be the set of all finite functions $\beta : M \rightarrow M$ that consists of exactly k distinct n -cycles. Then:

$$\begin{aligned}
A_n &= \{\alpha \in \text{Aut}(M) \mid \alpha \text{ has infinitely many orbits of size } n\} \\
&= \bigcap_{k=1}^{\infty} \{\alpha \in \text{Aut}(M) \mid \alpha \text{ has at least } k \text{ orbits of size } n\} \\
&= \bigcap_{k=1}^{\infty} \bigcup_{\beta \in B_{n,k}} [\beta]_{\text{Aut}(M)},
\end{aligned}$$

where indeed, $\bigcup_{\beta \in B_{n,k}} [\beta]_{\text{Aut}(M)}$ is dense in $\text{Aut}(M)$: take any finite $\gamma : M \rightarrow M$ such that $[\gamma]_{\text{Aut}(M)}$ is nonempty. Then also $\bigcup_{\beta \in B_{n,k}} [\beta]_{\text{Aut}(M)} \cap [\gamma]_{\text{Aut}(M)} \neq \emptyset$, one can easily construct a permutation that extends γ and have at least k many n -cycles.

Now we see that $A = \bigcap_{n=1}^{\infty} A_n$ is a comeagre set consisting of all functions that have infinitely many n -cycles for each n . The only thing left to show is that the set of functions with no infinite cycle is also comeagre. Indeed, for $m \in M$ let $B_m = \{\alpha \in \text{Aut}(M) \mid m \text{ has finite orbit in } \alpha\}$. This is an open dense set. It is a sum over basic open sets generated

by finite permutations with m in their domain. Denseness is also easy to see.

Finally, by the proposition 4.1, we can say that

$$\sigma^{\text{Aut}(M)} = \bigcap_{n=1}^{\infty} A_n \cap \bigcap_{m \in M} B_m,$$

which concludes the proof. \square

4.2. More general structures.

Fact 4.3. *Suppose M is an arbitrary structure and $f_1, f_2 \in \text{Aut}(M)$. Then f_1 and f_2 are conjugate if and only if $(M, f_1) \cong (M, f_2)$ as structures with one additional unary relation that is an automorphism.*

Proof. Suppose that $f_1 = g^{-1}f_2g$ for some $g \in \text{Aut}(M)$. Then g is the automorphism we're looking for. On the other hand if $g: (M, f_1) \rightarrow (M, f_2)$ is an isomorphism, then $g \circ f_1 = f_2 \circ g$ which exactly means that f_1, f_2 conjugate. \square

Theorem 4.4. *Let \mathcal{C} be a Fraïssé class of finite structures of a theory T in a relational language L . Let \mathcal{D} be the class of the finite structures of T in the language L with additional unary function symbol interpreted as an automorphism of the structure. If \mathcal{C} has the weak Hrushovski property and \mathcal{D} is a Fraïssé class, then there is a comeagre conjugacy class in the automorphism group of the $\text{Flim}(\mathcal{C})$.*

Proof. Let $\Gamma = \text{Flim}(\mathcal{C})$ and $(\Pi, \sigma) = \text{Flim}(\mathcal{D})$. Let $G = \text{Aut}(\Gamma)$, i.e. G is the automorphism group of Γ . First, by the theorem 3.21, we may assume without the loss of generality that $\Pi = \Gamma$. We will construct a strategy for the second player in the Banach-Mazur game on the topological space G . This strategy will give us a subset $A \subseteq G$ and as we will see a subset of a conjugacy class in G . By the Banach-Mazur theorem 2.11 this will prove that this class is comeagre.

Recall, G has a basis consisting of open sets $\{g \in G \mid g \upharpoonright_A = g_0 \upharpoonright_A\}$ for some finite set $A \subseteq \Gamma$ and some automorphism $g_0 \in G$. In other words, a basic open set is a set of all extensions of some finite partial isomorphism g_0 of Γ . By $B_g \subseteq G$ we denote a basic open subset given by a finite partial isomorphism g . From now on we will consider only finite partial isomorphism g such that B_g is nonempty.

With the use of corollary 2.15 we can consider only games, where both players choose finite partial isomorphisms. Namely, player I picks functions f_0, f_1, \dots and player II chooses g_0, g_1, \dots such that $f_0 \subseteq g_0 \subseteq f_1 \subseteq g_1 \subseteq \dots$, which identify the corresponding basic open subsets $B_{f_0} \supseteq B_{g_0} \supseteq \dots$

Our goal is to choose g_i in such a manner that the resulting function $g = \bigcap_{i=0}^{\infty} g_i$ will be an automorphism of the Fraïssé limit Γ such that $(\Gamma, g) \in \text{Flim} \mathcal{D}$. Precisely, $\bigcap_{i=0}^{\infty} B_{g_i} = \{g\}$, by the Fraïssé theorem 3.8

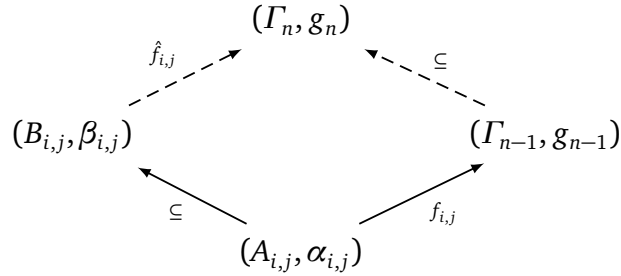
it will follow that $(\Gamma, g) \cong (\Pi, \sigma)$. Hence, by the fact 4.3, g and σ conjugate in G .

Once again, by the Fraïssé theorem 3.8 and the 3.10 lemma constructing g_i 's in a way such that age of (Γ, g) is exactly \mathcal{D} and so that it is weakly ultrahomogeneous will produce expected result. First, let us enumerate all pairs of structures $\{(A_n, \alpha_n), (B_n, \beta_n)\}_{n \in \mathbb{N}} \in \mathcal{D}$ such that the first element of the pair embeds by inclusion in the second, i.e. $(A_n, \alpha_n) \subseteq (B_n, \beta_n)$. Also, it may be that A_n is an empty. We enumerate the elements of the Fraïssé limit $\Gamma = \{v_0, v_1, \dots\}$.

Fix a bijection $\gamma: \mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N}$ such that for any $n, m \in \mathbb{N}$ we have $\gamma(n, m) \geq n$. This bijection naturally induces a well ordering on $\mathbb{N} \times \mathbb{N}$. This will prove useful later, as the main argument of the proof will be constructed as a bookkeeping argument.

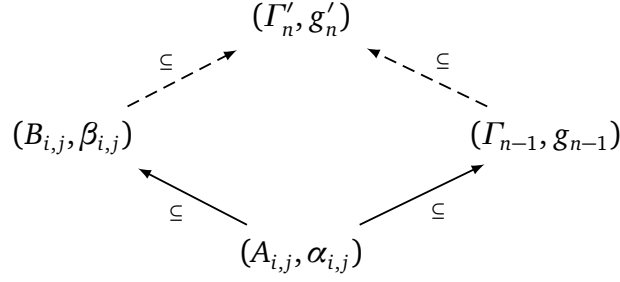
Just for sake of fixing a technical problem, let $g_{-1} = \emptyset$ and $X_{-1} = \emptyset$. Suppose that player I in the n -th move chooses a finite partial isomorphism f_n . We will construct $g_n \supseteq f_n$ and a set $X_n \subseteq \mathbb{N}^2$ such that following properties hold:

- (i) g_n is an automorphism of the induced substructure Γ_n ,
- (ii) $g_n(v_n)$ and $g_n^{-1}(v_n)$ are defined,
- (iii) let $\{(A_{n,k}, \alpha_{n,k}), (B_{n,k}, \beta_{n,k}), f_{n,k}\}_{k \in \mathbb{N}}$ be the enumeration of all pairs of finite structures of T with automorphism such that the first is a substructure of the second, i.e. $(A_{n,k}, \alpha_{n,k}) \subseteq (B_{n,k}, \beta_{n,k})$, and $f_{n,k}$ is an embedding of $(A_{n,k}, \alpha_{n,k})$ in the Γ_{n-1} (which is the substructure induced by g_{n-1}). Let $(i, j) = \min\{(\{0, 1, \dots\} \times \mathbb{N}) \setminus X_{n-1}\}$ (with the order induced by γ). Then $X_n = X_{n-1} \cup \{(i, j)\}$ and $(B_{n,k}, \beta_{n,k})$ embeds in (Γ_n, g_n) so that this diagram commutes:

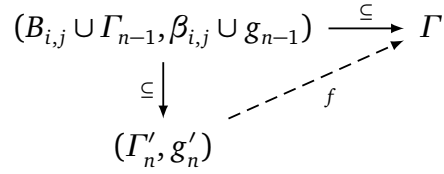


First item makes sure that no infinite orbit will be present in g . The second item together with the first one are necessary for g to be an automorphism of Γ . The third item is the one that gives weak ultrahomogeneity. Now we will show that indeed such g_n may be constructed.

First, we will suffice the item (iii). Namely, we will construct Γ'_n, g'_n such that $g_{n-1} \subseteq g'_n$ and $f_{i,j}$ extends to an embedding of $(B_{i,j}, \beta_{i,j})$ to (Γ'_n, g'_n) . But this can be easily done by the fact, that \mathcal{D} has the amalgamation property. Moreover, without the loss of generality we can assume that all embeddings are inclusions.



By the weak ultrahomogeneity we may assume that $\Gamma'_n \subseteq \Gamma$:



Now, by the WHP of \mathcal{H} we can extend the graph $\Gamma'_n \cup \{v_n\}$ together with its partial isomorphism g'_n to a graph Γ_n together with its automorphism $g_n \supseteq g'_n$ and without the loss of generality we may assume that $\Gamma_n \subseteq \Gamma$. This way we've constructed g_n that has all desired properties.

Now we need to see that $g = \bigcap_{n=0}^{\infty} g_n$ is indeed an automorphism of Γ such that (Γ, g) has the age \mathcal{H} and is weakly ultrahomogeneous. It is of course an automorphism of Γ as it is defined for every $v \in \Gamma$ and is a sum of increasing chain of finite automorphisms.

Take any finite structure of T with automorphism (B, β) . Then, there are i, j such that $(B, \beta) = (B_{i,j}, \beta_{i,j})$ and $A_{i,j} = \emptyset$. By the bookkeeping there was n such that $(i, j) = \min\{\{0, 1, \dots\} \times \mathbb{N} \setminus X_{n-1}\}$. This means that (B, β) embeds into (Γ_n, g_n) , hence it embeds into (Γ, g) , thus it has age \mathcal{H} . With a similar argument we can see that (Γ, g) is weakly ultrahomogeneous.

By this we know that g and σ conjugate in G . As we stated in the beginning of the proof, the set A of possible outcomes of the game (i.e. possible g 's we end up with) is comeagre in G , thus σ^G is also comeagre and σ is a generic automorphism, as it contains a comeagre set A . \square

4.3. Properties of the generic automorphism. Let \mathcal{C} be a Fraïssé class in a finite relational language L with weak Hrushovski property. Let \mathcal{H} be the Fraïssé class of the L -structures with additional automorphism symbol. Let $\Gamma = \text{Flim}(\mathcal{C})$.

Proposition 4.5. *Let σ be the generic automorphism of Γ . Then the set of fixed points of σ is isomorphic to Γ .*

Proof. Let $S = \{x \in \Gamma \mid \sigma(x) = x\}$. First we need to show that it is an infinite. By the theorem 4.4 we know that (Γ, σ) is the Fraïssé limit of \mathcal{H} , thus we can embed finite L -structures of any size with identity as an automorphism of the structure into (Γ, σ) . Thus S has to be infinite.

Also, the same argument shows that the age of the structure is exactly \mathcal{C} . It is weakly ultrahomogeneous, also by the fact that (Γ, σ) is in \mathcal{H} . \square

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